Part VI

æ.,

イロト イヨト イヨト イヨト

History of Probabilistic Numerical Methods



Bayesian Solution of Ordinary Differential Equations

J. SKILLING, Cambridge University.

In: Maximum Entropy and Bayesian Methods, Springer Netherlands, 23-37, 1992.

In the numerical solution of ordinary differential equations, a function y(x) is to be reconstructed from knowledge of the functional form of its derivative: dy/dx = f(x, y), together with an appropriate boundary condition. The derivative f is evaluated at a sequence of suitably chosen points (x_k, y_k) , from which the form of $y(\cdot)$ is estimated. This is an inference problem, which can and perhaps should be treated by Bayesian techniques. As always, the inference appears as a probability distribution $prob(y(\cdot))$, from which random samples show the probabilistic reliability of the results.

History of Probabilistic Numerical Methods



Bayesian Solution of Ordinary Differential Equations

J. SKILLING, Cambridge University.

In: Maximum Entropy and Bayesian Methods, Springer Netherlands, 23-37, 1992.

In the numerical solution of ordinary differential equations, a function y(x) is to be reconstructed from knowledge of the functional form of its derivative: dy/dx = f(x, y), together with an appropriate boundary condition. The derivative f is evaluated at a sequence of suitably chosen points (x_k, y_k) , from which the form of $y(\cdot)$ is estimated. This is an inference problem, which can and perhaps should be treated by Bayesian techniques. As always, the inference appears as a probability distribution $prob(y(\cdot))$, from which random samples show the probabilistic reliability of the results.

Twelfth Job: Introduction to Graphical Models

A B > 4
 B > 4
 B

Numerical analysis for the "drag and drop" era of computational pipelines:



[Fig: IBM High Performance Computation]

< □ > < ^[] >

The sophistication and scale of modern computer models creates an urgent need to better understand the propagation and accumulation of numerical error within arbitrary - often large - pipelines of computation, so that "numerical risk" to end-users can be controlled.

→ Need to consider graphical representations of computation.

Numerical analysis for the "drag and drop" era of computational pipelines:



[Fig: IBM High Performance Computation]

The sophistication and scale of modern computer models creates an urgent need to better understand the propagation and accumulation of numerical error within arbitrary - often large - pipelines of computation, so that "numerical risk" to end-users can be controlled.

 \implies Need to consider graphical representations of computation.

• A graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, with vertex set $\mathcal{V} = \{v_1, \dots, v_p\}$ and edge set $\mathcal{E} \subseteq \mathcal{V} \times \mathcal{V}$.

 Semantics that associate the graph to statements about conditional (in)dependence of random variables {X_v}_{v∈V}.

Recall: X is C.I. of Y given $Z \leftrightarrow X \perp \!\!\!\perp Y | Z \leftrightarrow p_{X,Y|Z}(x,y) = p_{X|Z}(x)p_{Y|Z}(y)$

Example: In a Gaussian graphical model:

- The random variables $\{X_v\}_{v \in \mathcal{V}}$ are jointly Gaussian.
- The edges are undirected (i.e. $(i, j) \in \mathcal{E}$ iff $(j, i) \in \mathcal{E}$).
- Two vertices v_1 and v_2 are **not** connected by an edge iff $X_{v_1} \perp \!\!\!\perp X_{v_2} | X_{\mathcal{V} \setminus \{v_1, v_2\}}$.

Fact: The edge structure of a Gaussian graphical model characterises the sparsity structure of the associated precision matrix.

• A graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, with vertex set $\mathcal{V} = \{v_1, \dots, v_p\}$ and edge set $\mathcal{E} \subseteq \mathcal{V} \times \mathcal{V}$.

 Semantics that associate the graph to statements about conditional (in)dependence of random variables {X_v}_{v∈V}.

Recall: X is C.I. of Y given $Z \leftrightarrow X \perp Y | Z \leftrightarrow p_{X,Y|Z}(x,y) = p_{X|Z}(x)p_{Y|Z}(y)$

Example: In a Gaussian graphical model:

- The random variables $\{X_v\}_{v \in \mathcal{V}}$ are jointly Gaussian.
- The edges are undirected (i.e. $(i,j) \in \mathcal{E}$ iff $(j,i) \in \mathcal{E}$).
- Two vertices v_1 and v_2 are **not** connected by an edge iff $X_{v_1} \perp \!\!\!\perp X_{v_2} | X_{\mathcal{V} \setminus \{v_1, v_2\}}$.

Fact: The edge structure of a Gaussian graphical model characterises the sparsity structure of the associated precision matrix.

• A graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, with vertex set $\mathcal{V} = \{v_1, \dots, v_p\}$ and edge set $\mathcal{E} \subseteq \mathcal{V} \times \mathcal{V}$.

 Semantics that associate the graph to statements about conditional (in)dependence of random variables {X_v}_{v∈V}.

Recall: X is C.I. of Y given $Z \leftrightarrow X \perp Y | Z \leftrightarrow p_{X,Y|Z}(x,y) = p_{X|Z}(x)p_{Y|Z}(y)$

Example: In a Gaussian graphical model:

- The random variables $\{X_v\}_{v \in \mathcal{V}}$ are jointly Gaussian.
- The edges are undirected (i.e. $(i,j) \in \mathcal{E}$ iff $(j,i) \in \mathcal{E}$).
- Two vertices v_1 and v_2 are **not** connected by an edge iff $X_{v_1} \perp \!\!\!\perp X_{v_2} | X_{\mathcal{V} \setminus \{v_1, v_2\}}$.

Fact: The edge structure of a Gaussian graphical model characterises the sparsity structure of the associated precision matrix.

• A graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, with vertex set $\mathcal{V} = \{v_1, \dots, v_p\}$ and edge set $\mathcal{E} \subseteq \mathcal{V} \times \mathcal{V}$.

 Semantics that associate the graph to statements about conditional (in)dependence of random variables {X_v}_{v∈V}.

Recall: X is C.I. of Y given $Z \leftrightarrow X \perp \!\!\!\perp Y | Z \leftrightarrow p_{X,Y|Z}(x,y) = p_{X|Z}(x)p_{Y|Z}(y)$

Example: In a Gaussian graphical model:

- The random variables $\{X_v\}_{v \in \mathcal{V}}$ are jointly Gaussian.
- The edges are undirected (i.e. $(i,j) \in \mathcal{E}$ iff $(j,i) \in \mathcal{E}$).
- Two vertices v_1 and v_2 are **not** connected by an edge iff $X_{v_1} \perp \!\!\!\perp X_{v_2} | X_{\mathcal{V} \setminus \{v_1, v_2\}}$.

Fact: The edge structure of a Gaussian graphical model characterises the sparsity structure of the associated precision matrix.

• A graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, with vertex set $\mathcal{V} = \{v_1, \dots, v_p\}$ and edge set $\mathcal{E} \subseteq \mathcal{V} \times \mathcal{V}$.

 Semantics that associate the graph to statements about conditional (in)dependence of random variables {X_v}_{v∈V}.

NB: More complicated graphical structures have also been developed.

Example: In a Gaussian graphical model:

- The random variables $\{X_v\}_{v \in \mathcal{V}}$ are jointly Gaussian.
- The edges are undirected (i.e. $(i,j) \in \mathcal{E}$ iff $(j,i) \in \mathcal{E}$).
- Two vertices v_1 and v_2 are **not** connected by an edge iff $X_{v_1} \perp \!\!\!\perp X_{v_2} | X_{\mathcal{V} \setminus \{v_1, v_2\}}$.

Fact: The edge structure of a Gaussian graphical model characterises the sparsity structure of the associated precision matrix.

• A graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, with vertex set $\mathcal{V} = \{v_1, \dots, v_p\}$ and edge set $\mathcal{E} \subseteq \mathcal{V} \times \mathcal{V}$.

 Semantics that associate the graph to statements about conditional (in)dependence of random variables {X_v}_{v∈V}.

NB: More complicated graphical structures have also been developed.

Example: In a Gaussian graphical model:

- The random variables $\{X_v\}_{v \in \mathcal{V}}$ are jointly Gaussian.
- The edges are undirected (i.e. $(i, j) \in \mathcal{E}$ iff $(j, i) \in \mathcal{E}$).
- Two vertices v_1 and v_2 are **not** connected by an edge iff $X_{v_1} \perp \!\!\!\perp X_{v_2} | X_{\mathcal{V} \setminus \{v_1, v_2\}}$.

Fact: The edge structure of a Gaussian graphical model characterises the sparsity structure of the associated precision matrix.

• A graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, with vertex set $\mathcal{V} = \{v_1, \dots, v_p\}$ and edge set $\mathcal{E} \subseteq \mathcal{V} \times \mathcal{V}$.

 Semantics that associate the graph to statements about conditional (in)dependence of random variables {X_v}_{v∈V}.

NB: More complicated graphical structures have also been developed.

Example: In a Gaussian graphical model:

- The random variables $\{X_{\nu}\}_{\nu \in \mathcal{V}}$ are jointly Gaussian.
- The edges are undirected (i.e. $(i, j) \in \mathcal{E}$ iff $(j, i) \in \mathcal{E}$).
- Two vertices v_1 and v_2 are **not** connected by an edge iff $X_{v_1} \perp \!\!\!\perp X_{v_2} | X_{\mathcal{V} \setminus \{v_1, v_2\}}$.

Fact: The edge structure of a Gaussian graphical model characterises the sparsity structure of the associated precision matrix.

- The edges are directed (i.e. $(i,j) \in \mathcal{E} \implies (j,i) \notin \mathcal{E}$).
- No directed cycles exist (in particular, $(i, i) \notin \mathcal{E}$).
- Two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff $X_Q \perp \!\!\!\perp X_R | X_S$. ("global Markov property")

Let $\operatorname{an}_{\mathcal{G}}(S)$ denote the *ancestors* of the set S in \mathcal{G} . (NB: This includes S itself.)

Given a directed, acyclic graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, we say that two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff the following holds:

- Form \mathcal{G}' , the subgraph induced by $\operatorname{an}_{\mathcal{G}}(Q \cup R \cup S)$.
- Form \mathcal{G}'' by connecting any un-connected parents (with an undirected edge) in \mathcal{G}' .
- Remove all arrowheads in \mathcal{G}'' , to obtain an undirected "moral" graph \mathcal{G}_{moral} .
- Check that all paths from vertices in *Q* to vertices in *R* in *G*_{moral} pass through a vertex in *S*.

• • • • • • • • • • • •

- The edges are directed (i.e. $(i,j) \in \mathcal{E} \implies (j,i) \notin \mathcal{E}$).
- No directed cycles exist (in particular, $(i, i) \notin \mathcal{E}$).
- Two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff $X_Q \perp \!\!\!\perp X_R | X_S$. ("global Markov property")

Let $\operatorname{an}_{\mathcal{G}}(S)$ denote the *ancestors* of the set S in \mathcal{G} . (NB: This includes S itself.)

Given a directed, acyclic graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, we say that two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff the following holds:

- Form \mathcal{G}' , the subgraph induced by $\operatorname{an}_{\mathcal{G}}(Q \cup R \cup S)$.
- Form \mathcal{G}'' by connecting any un-connected parents (with an undirected edge) in \mathcal{G}' .
- Remove all arrowheads in \mathcal{G}'' , to obtain an undirected "moral" graph \mathcal{G}_{moral} .
- Check that all paths from vertices in *Q* to vertices in *R* in *G*_{moral} pass through a vertex in *S*.

• • • • • • • • • • • •

- The edges are directed (i.e. $(i,j) \in \mathcal{E} \implies (j,i) \notin \mathcal{E}$).
- No directed cycles exist (in particular, $(i, i) \notin \mathcal{E}$).
- Two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff $X_Q \perp \!\!\!\perp X_R | X_S$. ("global Markov property")

Let $\operatorname{an}_{\mathcal{G}}(S)$ denote the *ancestors* of the set S in \mathcal{G} . (NB: This includes S itself.)

Given a directed, acyclic graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, we say that two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff the following holds:

- Form \mathcal{G}' , the subgraph induced by $\operatorname{an}_{\mathcal{G}}(Q \cup R \cup S)$.
- Form \mathcal{G}'' by connecting any un-connected parents (with an undirected edge) in \mathcal{G}' .
- Remove all arrowheads in \mathcal{G}'' , to obtain an undirected "moral" graph \mathcal{G}_{moral} .
- Check that all paths from vertices in *Q* to vertices in *R* in *G*_{moral} pass through a vertex in *S*.

- The edges are directed (i.e. $(i,j) \in \mathcal{E} \implies (j,i) \notin \mathcal{E}$).
- No directed cycles exist (in particular, $(i, i) \notin \mathcal{E}$).
- Two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff $X_Q \perp \!\!\!\perp X_R | X_S$. ("global Markov property")

Let $\operatorname{an}_{\mathcal{G}}(S)$ denote the *ancestors* of the set S in \mathcal{G} . (NB: This includes S itself.)

Given a directed, acyclic graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, we say that two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff the following holds:

- Form \mathcal{G}' , the subgraph induced by $\operatorname{an}_{\mathcal{G}}(Q \cup R \cup S)$.
- Form \mathcal{G}'' by connecting any un-connected parents (with an undirected edge) in \mathcal{G}' .
- ullet Remove all arrowheads in ${\cal G}''$, to obtain an undirected "moral" graph ${\cal G}_{\sf moral}$.
- Check that all paths from vertices in Q to vertices in R in $\mathcal{G}_{\text{moral}}$ pass through a vertex in S.

- The edges are directed (i.e. $(i,j) \in \mathcal{E} \implies (j,i) \notin \mathcal{E}$).
- No directed cycles exist (in particular, $(i, i) \notin \mathcal{E}$).
- Two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff $X_Q \perp \!\!\!\perp X_R | X_S$. ("global Markov property")

Let $\operatorname{an}_{\mathcal{G}}(S)$ denote the *ancestors* of the set S in \mathcal{G} . (NB: This includes S itself.)

Given a directed, acyclic graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, we say that two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff the following holds:

- Form \mathcal{G}' , the subgraph induced by $\operatorname{an}_{\mathcal{G}}(Q \cup R \cup S)$.
- Form \mathcal{G}'' by connecting any un-connected parents (with an undirected edge) in \mathcal{G}' .
- ullet Remove all arrowheads in \mathcal{G}'' , to obtain an undirected "moral" graph $\mathcal{G}_{\mathsf{moral}}$
- Check that all paths from vertices in Q to vertices in R in $\mathcal{G}_{\text{moral}}$ pass through a vertex in S.

- The edges are directed (i.e. $(i,j) \in \mathcal{E} \implies (j,i) \notin \mathcal{E}$).
- No directed cycles exist (in particular, $(i, i) \notin \mathcal{E}$).
- Two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff $X_Q \perp \!\!\!\perp X_R | X_S$. ("global Markov property")

Let $\operatorname{an}_{\mathcal{G}}(S)$ denote the *ancestors* of the set S in \mathcal{G} . (NB: This includes S itself.)

Given a directed, acyclic graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, we say that two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff the following holds:

- Form \mathcal{G}' , the subgraph induced by $\operatorname{an}_{\mathcal{G}}(Q \cup R \cup S)$.
- Form \mathcal{G}'' by connecting any un-connected parents (with an undirected edge) in \mathcal{G}' .
- Remove all arrowheads in \mathcal{G}'' , to obtain an undirected "moral" graph $\mathcal{G}_{\mathsf{moral}}$.
- Check that all paths from vertices in Q to vertices in R in $\mathcal{G}_{\text{moral}}$ pass through a vertex in S.

- The edges are directed (i.e. $(i,j) \in \mathcal{E} \implies (j,i) \notin \mathcal{E}$).
- No directed cycles exist (in particular, $(i, i) \notin \mathcal{E}$).
- Two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff $X_Q \perp \!\!\!\perp X_R | X_S$. ("global Markov property")

Let $\operatorname{an}_{\mathcal{G}}(S)$ denote the *ancestors* of the set S in \mathcal{G} . (NB: This includes S itself.)

Given a directed, acyclic graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, we say that two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff the following holds:

- Form \mathcal{G}' , the subgraph induced by $\operatorname{an}_{\mathcal{G}}(Q \cup R \cup S)$.
- Form \mathcal{G}'' by connecting any un-connected parents (with an undirected edge) in \mathcal{G}' .
- \bullet Remove all arrowheads in $\mathcal{G}'',$ to obtain an undirected "moral" graph $\mathcal{G}_{\text{moral}}.$
- Check that all paths from vertices in Q to vertices in R in $\mathcal{G}_{\text{moral}}$ pass through a vertex in S.

- The edges are directed (i.e. $(i,j) \in \mathcal{E} \implies (j,i) \notin \mathcal{E}$).
- No directed cycles exist (in particular, $(i, i) \notin \mathcal{E}$).
- Two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff $X_Q \perp \!\!\!\perp X_R | X_S$. ("global Markov property")

Let $\operatorname{an}_{\mathcal{G}}(S)$ denote the *ancestors* of the set S in \mathcal{G} . (NB: This includes S itself.)

Given a directed, acyclic graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, we say that two sets $Q, R \subseteq \mathcal{V}$ are d-separated by a set $S \subseteq \mathcal{V} \setminus Q \cup R$ iff the following holds:

- Form \mathcal{G}' , the subgraph induced by $\operatorname{an}_{\mathcal{G}}(Q \cup R \cup S)$.
- Form \mathcal{G}'' by connecting any un-connected parents (with an undirected edge) in \mathcal{G}' .
- \bullet Remove all arrowheads in \mathcal{G}'' , to obtain an undirected "moral" graph $\mathcal{G}_{moral}.$
- Check that all paths from vertices in Q to vertices in R in \mathcal{G}_{moral} pass through a vertex in S.

Exercise: Is Vent-sys d-separated from HTPS?



2

Local Markov Property

For a Bayesian network;

$$X_v \perp \!\!\!\perp X_{\operatorname{Nd}(v) \setminus \operatorname{Pa}(v)} | X_{\operatorname{Pa}(v)}$$

where, according to the graph \mathcal{G} ,

- Nd(v) are the non-descendants of $v \in V$
- Pa(v) are the parents of $v \in \mathcal{V}$.

Message: "Only immediate parents matter".

Practically important, as it allows local verification of the global Markov property.

Local Markov Property

For a Bayesian network;

$$X_v \perp \!\!\!\perp X_{\operatorname{Nd}(v) \setminus \operatorname{Pa}(v)} | X_{\operatorname{Pa}(v)}$$

where, according to the graph \mathcal{G} ,

- Nd(v) are the non-descendants of $v \in V$
- Pa(v) are the parents of $v \in \mathcal{V}$.

Message: "Only immediate parents matter".

Practically important, as it allows local verification of the global Markov property.

Local Markov Property

For a Bayesian network;

$$X_{v} \perp \perp X_{\operatorname{Nd}(v) \setminus \operatorname{Pa}(v)} | X_{\operatorname{Pa}(v)}$$

where, according to the graph \mathcal{G} ,

- Nd(v) are the non-descendants of $v \in \mathcal{V}$
- Pa(v) are the parents of $v \in \mathcal{V}$.

Message: "Only immediate parents matter".

Practically important, as it allows local verification of the global Markov property.

Is there a connection to computation?



[Image from Li Haoyi's blog]

In particular, is there a connection to functional programming?



[Image from Li Haoyi's blog]

Or a connection to grid or cloud computing?



Thirteenth Job: Pipelines of Computation

A B > 4
 B > 4
 B

Consider estimation of the integral

with a Bayesian probabilistic numerical method (Bayesian quadrature), based on the information $\{x(t_1), \ldots, x(t_{2m})\}$, where $t_1 = 0, t_m = 0.5, t_{2m} = 1$.

 $\int_{a}^{1} x(t) \mathrm{d}t$

Under what circumstances can this computation be partitioned into two independent sub-computations?

$$\int_0^1 u(x) \mathrm{d}x = \int_0^{0.5} u(x) \mathrm{d}x + \int_{0.5}^1 u(x) \mathrm{d}x$$

Consider estimation of the integral

 $\int_0^1 x(t) \mathrm{d}t$

with a Bayesian probabilistic numerical method (Bayesian quadrature), based on the information $\{x(t_1), \ldots, x(t_{2m})\}$, where $t_1 = 0, t_m = 0.5, t_{2m} = 1$.

Under what circumstances can this computation be partitioned into two independent sub-computations?

$$\int_0^1 u(x) \mathrm{d}x = \int_0^{0.5} u(x) \mathrm{d}x + \int_{0.5}^1 u(x) \mathrm{d}x$$

Let's attempt to represent this with a graphical model, which we will call a pipeline:



- Nodes are of two kinds: *Information* nodes □, and *method* nodes ■.
- The graph is bipartite, so that edges connect a method node to an information node or vice-versa. That is, edges are of the form □ → ■ or ■ → □.
- There are in general *n* method nodes, each with a unique label in $\{1, \ldots, n\}$.
- The method node labelled *i* has *m*(*i*) parents and one child. Its in-edges are assigned a unique label in {1,..., *m*(*i*)}.
- There is a unique terminal node and it is the child of method node *n*. This represents the principal quantity of interest $Q(x) = \int_0^1 x(t) dt$.

A = A = A = A = A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A

Let's attempt to represent this with a graphical model, which we will call a pipeline:



Nodes are of two kinds: Information nodes □, and method nodes ■.

- The graph is bipartite, so that edges connect a method node to an information node or vice-versa. That is, edges are of the form □ → ■ or ■ → □.
- There are in general *n* method nodes, each with a unique label in $\{1, \ldots, n\}$.
- The method node labelled *i* has *m*(*i*) parents and one child. Its in-edges are assigned a unique label in {1,..., *m*(*i*)}.
- There is a unique terminal node and it is the child of method node *n*. This represents the principal quantity of interest $Q(x) = \int_0^1 x(t) dt$.

A = A = A = A = A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A

Let's attempt to represent this with a graphical model, which we will call a pipeline:



- Nodes are of two kinds: Information nodes □, and method nodes ■.
- The graph is bipartite, so that edges connect a method node to an information node or vice-versa. That is, edges are of the form □ → ■ or ■ → □.
- There are in general *n* method nodes, each with a unique label in $\{1, \ldots, n\}$.
- The method node labelled *i* has *m*(*i*) parents and one child. Its in-edges are assigned a unique label in {1,..., *m*(*i*)}.
- There is a unique terminal node and it is the child of method node *n*. This represents the principal quantity of interest $Q(x) = \int_0^1 x(t) dt$.

A B > A B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A
 B > A

Let's attempt to represent this with a graphical model, which we will call a pipeline:



- Nodes are of two kinds: *Information* nodes □, and *method* nodes ■.
- The graph is bipartite, so that edges connect a method node to an information node or vice-versa. That is, edges are of the form □ → ■ or ■ → □.
- There are in general n method nodes, each with a unique label in $\{1, \ldots, n\}$.
- The method node labelled *i* has *m*(*i*) parents and one child. Its in-edges are assigned a unique label in {1,..., *m*(*i*)}.
- There is a unique terminal node and it is the child of method node *n*. This represents the principal quantity of interest $Q(x) = \int_0^1 x(t) dt$.

A = A = A = A = A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A
 A

Let's attempt to represent this with a graphical model, which we will call a pipeline:



- Nodes are of two kinds: *Information* nodes □, and *method* nodes ■.
- The graph is bipartite, so that edges connect a method node to an information node or vice-versa. That is, edges are of the form □ → ■ or ■ → □.
- There are in general *n* method nodes, each with a unique label in $\{1, \ldots, n\}$.
- The method node labelled *i* has *m*(*i*) parents and one child. Its in-edges are assigned a unique label in {1,..., *m*(*i*)}.
- There is a unique terminal node and it is the child of method node *n*. This represents the principal quantity of interest $Q(x) = \int_0^1 x(t) dt$.

A D > A P > A B > A

Let's attempt to represent this with a graphical model, which we will call a pipeline:



- Nodes are of two kinds: *Information* nodes □, and *method* nodes ■.
- The graph is bipartite, so that edges connect a method node to an information node or vice-versa. That is, edges are of the form □ → ■ or ■ → □.
- There are in general n method nodes, each with a unique label in $\{1, \ldots, n\}$.
- The method node labelled *i* has *m*(*i*) parents and one child. Its in-edges are assigned a unique label in {1,..., *m*(*i*)}.
- There is a unique terminal node and it is the child of method node *n*. This represents the principal quantity of interest $Q(x) = \int_0^1 x(t) dt$.

Let's attempt to represent this with a graphical model, which we will call a pipeline:



Let

- M_1 be Bayesian quadrature for $\int_0^{0.5} x(t) dt$.
- M_2 be Bayesian quadrature for $\int_{0.5}^{1} x(t) dt$.
- M_3 be the trivial probabilistic numerical method that adds its two inputs..

Let's attempt to represent this with a graphical model, which we will call a pipeline:



Q: When is the output of the pipeline Bayesian?

I.e. When does the output of the pipeline coincide with standard Bayesian quadrature performed on the full information $\{x(t_1), \ldots, x(t_{2m})\}$?

< □ > < ^[] >



If we restrict attention to Bayesian probabilistic numerical methods, then M_1 , M_2 and M_3 are uniquely determined by the prior distribution P_x for the integrand.

So we can delete the method nodes to obtain the dependency graph of a pipeline.

Starting to look like a Bayesian network ...

A B A B A
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 B
 A
 A
 A
 A
 A
 A



If we restrict attention to Bayesian probabilistic numerical methods, then M_1 , M_2 and M_3 are uniquely determined by the prior distribution P_x for the integrand.

So we can delete the method nodes to obtain the dependency graph of a pipeline.

Starting to look like a Bayesian network ...

A D > A B > A B >



If we restrict attention to Bayesian probabilistic numerical methods, then M_1 , M_2 and M_3 are uniquely determined by the prior distribution P_x for the integrand.

So we can delete the method nodes to obtain the dependency graph of a pipeline.

Starting to look like a Bayesian network ...

A D > A B > A B >



If we restrict attention to Bayesian probabilistic numerical methods, then M_1 , M_2 and M_3 are uniquely determined by the prior distribution P_x for the integrand.

So we can delete the method nodes to obtain the dependency graph of a pipeline.

Starting to look like a Bayesian network...



Associate each node i with a random variable X_i .

e.g.
$$X_1 = \{x(t_1), \dots, x(t_{m-1})\}, X_4 = \int_0^{0.5} x(t) dt.$$

A prior P_x is <u>coherent</u> for the dependency graph if each $X_i \perp X_{Nd(i) \setminus Pa(i)} \mid X_{Pa(i)}$.

• • • • • • • •



Associate each node i with a random variable X_i .

e.g.
$$X_1 = \{x(t_1), \dots, x(t_{m-1})\}, X_4 = \int_0^{0.5} x(t) dt.$$

A prior P_x is <u>coherent</u> for the dependency graph if each $X_i \perp X_{Nd(i) \setminus Pa(i)} \mid X_{Pa(i)}$.

A ID > A ID > A

Split Integration: Coherence



Thus in this instance we would ask whether $\int_{0.5}^{1} x(t) dt$ is independent of $x(t_1), \ldots, x(t_{m-1})$ given $x(t_m), \ldots, x(t_{2m})$?

This is <u>not</u> true in general, but sometimes holds - e.g. a Wiener process prior P_x .

Split Integration: Coherence



Thus in this instance we would ask whether $\int_{0.5}^{1} x(t) dt$ is independent of $x(t_1), \ldots, x(t_{m-1})$ given $x(t_m), \ldots, x(t_{2m})$?

This is <u>not</u> true in general, but sometimes holds - e.g. a Wiener process prior P_{x} .

< □ > < ^[] >

Split Integration: Coherence



Thus in this instance we would ask whether $\int_{0.5}^{1} x(t) dt$ is independent of $x(t_1), \ldots, x(t_{m-1})$ given $x(t_m), \ldots, x(t_{2m})$?

This is <u>not</u> true in general, but sometimes holds - e.g. a Wiener process prior P_{x} .

The process illustrated here can be made formal:

A pipeline is Bayesian for estimation of its output if:

- The prior P_x is <u>coherent</u> for the dependency graph associated to the pipeline.
- **(2)** The methods M_i are Bayesian probabilistic numerical methods.

Open Question: Can a similar notion of coherence be developed for non-Bayesian probabilistic numerical methods?

Point for Discussion: How to do split integration in $d \ge 2$ dimensions?

The process illustrated here can be made formal:

A pipeline is Bayesian for estimation of its output if:

- The prior P_x is <u>coherent</u> for the dependency graph associated to the pipeline.
- **2** The methods M_i are Bayesian probabilistic numerical methods.

Open Question: Can a similar notion of coherence be developed for non-Bayesian probabilistic numerical methods?

Point for Discussion: How to do split integration in $d \ge 2$ dimensions?

The process illustrated here can be made formal:

A pipeline is Bayesian for estimation of its output if:

- The prior P_x is <u>coherent</u> for the dependency graph associated to the pipeline.
- **(2)** The methods *M_i* are Bayesian probabilistic numerical methods.

Open Question: Can a similar notion of coherence be developed for non-Bayesian probabilistic numerical methods?

Point for Discussion: How to do split integration in $d \ge 2$ dimensions?

- Computational work-flow can be related to graphical models used in statistical applications.
- Bayesian probabilistic numerical methods induce a joint distribution over unknown objects, whose conditional (in)dependence structure can be represented with a pipeline graph.
- The local Markov property can be used to check whether a large pipeline of Bayesian probabilistic numerical methods is coherent.

END OF PART VI

- Computational work-flow can be related to graphical models used in statistical applications.
- Bayesian probabilistic numerical methods induce a joint distribution over unknown objects, whose conditional (in)dependence structure can be represented with a pipeline graph.
- The local Markov property can be used to check whether a large pipeline of Bayesian probabilistic numerical methods is coherent.

END OF PART VI

- Computational work-flow can be related to graphical models used in statistical applications.
- Bayesian probabilistic numerical methods induce a joint distribution over unknown objects, whose conditional (in)dependence structure can be represented with a pipeline graph.
- The local Markov property can be used to check whether a large pipeline of Bayesian probabilistic numerical methods is coherent.

END OF PART VI

Image: A match the second s

- Computational work-flow can be related to graphical models used in statistical applications.
- Bayesian probabilistic numerical methods induce a joint distribution over unknown objects, whose conditional (in)dependence structure can be represented with a pipeline graph.
- The local Markov property can be used to check whether a large pipeline of Bayesian probabilistic numerical methods is coherent.

END OF PART VI

Image: A match the second s

- Computational work-flow can be related to graphical models used in statistical applications.
- Bayesian probabilistic numerical methods induce a joint distribution over unknown objects, whose conditional (in)dependence structure can be represented with a pipeline graph.
- The local Markov property can be used to check whether a large pipeline of Bayesian probabilistic numerical methods is coherent.

END OF PART VI

Image: A match a ma

"All uncertainty is of one kind" \sim Phil Dawid.

Download link: oates.work/dobbiaco

< □ > < ^[] >

- J. Chang and D. Pollard. Conditioning as disintegration. *Statistica Neerlandica*, 51(3):287–317, 1997.
- I. Cialenco, G. E. Fasshauer, and Q. Ye. Approximation of stochastic partial differential equations by a kernel-based collocation method. *International Journal of Computer Mathematics*, 89 (18):2543–2561, 2012.
- J. Cockayne, C. Oates, T. J. Sullivan, and M. Girolami. Probabilistic meshless methods for partial differential equations and Bayesian inverse problems, 2016. arXiv:1605.07811v1.
- J. Cockayne, C. Oates, T. Sullivan, and M. Girolami. Bayesian probabilistic numerical methods. arXiv:1702.03673, 2017.
- P. R. Conrad, M. Girolami, S. Särkkä, A. M. Stuart, and K. C. Zygalakis. Statistical analysis of differential equations: introducing probability measures on numerical solutions. *Statistics and Computing*, 2016.
- S. L. Cotter, G. O. Roberts, A. M. Stuart, D. White, et al. Mcmc methods for functions: modifying old algorithms to make them faster. *Statistical Science*, 28(3):424–446, 2013.
- T. Cui, J. Martin, Y. M. Marzouk, A. Solonen, and A. Spantini. Likelihood-informed dimension reduction for nonlinear inverse problems. *Inverse Problems*, 30(11):114015, 2014.
- C. Dellacherie and P. Meyer. Probabilities and Potential. North-Holland, Amsterdam, 1978.
- T. E. Hull and J. R. Swenson. Tests of probabilistic models for propagation of roundoff errors. *Communications of the ACM*, 9(2):108–113, 1966.

イロト イポト イヨト イヨト

- D. Isaacson, J. L. Mueller, J. C. Newell, and S. Siltanen. Reconstructions of chest phantoms by the d-bar method for electrical impedance tomography. *IEEE Transactions on Medical Imaging*, 23(7):821–828, 2004.
- J. B. Kadane and G. W. Wasilkowski. *Bayesian Statistics*, chapter Average Case *ϵ*-Complexity in Computer Science: A Bayesian View, pages 361–374. Elsevier, North-Holland, 1985.
- A. N. Kolmogorov. Foundations of Probability. 1933.
- C. Ley, G. Reinert, Y. Swan, et al. Stein's method for comparison of univariate distributions. Probability Surveys, 14:1–52, 2017.
- C. Oates, S. Niederer, A. Lee, F. Briol, and M. Girolami. Probabilistic models for integration error in assessment of functional cardiac models. *arXiv:1606.06841*, 2017.
- C. J. Oates, J. Cockayne, and R. G. Aykroyd. Bayesian probabilistic numerical methods for industrial process monitoring. *In preparation*, 2016a.
- C. J. Oates, J. Cockayne, F.-X. Briol, and M. Girolami. Convergence rates for a class of estimators based on stein's identity. *arXiv:1603.03220*, 2016b.
- H. Owhadi. Bayesian numerical homogenization. *Multiscale Modeling & Simulation*, 13(3): 812–828, 2015.
- S. Särkkä. Linear operators and stochastic partial differential equations in gaussian process regression. Artificial Neural Networks and Machine Learning–ICANN 2011, pages 151–158, 2011.
- A. M. Stuart. Inverse problems: A Bayesian perspective. Acta Numerica, 19:451-559, May 2010.

イロト 不得下 イヨト イヨト

- T. J. Sullivan. Well-posed Bayesian inverse problems and heavy-tailed stable quasi-Banach space priors, 2016. arXiv:1605.05898.
- B. Szabó, A. van der Vaart, J. van Zanten, et al. Frequentist coverage of adaptive nonparametric bayesian credible sets. *The Annals of Statistics*, 43(4):1391–1428, 2015.
- H. Wendland. Scattered data approximation, volume 17. Cambridge university press, 2004.

• • • • • • • • • • • •